

Lower bounding the Folkman numbers

$$F_v(a_1, \dots, a_s; m - 1)$$

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Notations:

- $V(G)$ - the vertex set of G .
- $E(G)$ - the edge set of G .
- \overline{G} - the complement of G .
- K_n - complete graph on n vertices.
- C_n - simple cycle on n vertices.
- $G_1 + G_2$ - sum of the graphs G_1 and G_2 .

Definition. The expression $G \xrightarrow{v} (a_1, \dots, a_s)$ means that in every coloring of the vertices of the graph G in s colors (s -coloring) there exists $i \in \{1, \dots, s\}$ such that there is a monochromatic a_i -clique of color i .

- For example, $G \xrightarrow{v} (3, 3)$ means that in every coloring of the vertices of G in 2 colors there is a monochromatic triangle.

- $G \xrightarrow{v} \underbrace{(2, \dots, 2)}_r \iff \chi(G) > r$. Therefore, the problem of checking the property $G \xrightarrow{v} (a_1, \dots, a_s)$ can be considered a generalization of the problem of computing the chromatic number of G . For convenience, instead of $G \xrightarrow{v} \underbrace{(2, \dots, 2)}_r$ we shall write $G \xrightarrow{v} (2_r)$.

Definition. The vertex Folkman number $F_v(a_1, \dots, a_s; q)$ is defined by

$$F_v(a_1, \dots, a_s; q) = \min\{|V(G)| : G \xrightarrow{v} (a_1, \dots, a_s) \text{ and } K_q \not\subseteq G\}.$$

1970 Folkman:

$$F_v(a_1, \dots, a_s; q) \text{ exists} \Leftrightarrow q > \max\{a_1, \dots, a_s\}.$$

Denote:

$$\mathcal{H}_v(a_1, \dots, a_s; q) = \{G : G \xrightarrow{v} (a_1, \dots, a_s) \text{ and } K_q \not\subseteq G\}.$$

$$\mathcal{H}_v(a_1, \dots, a_s; q; n) = \{G : G \in \mathcal{H}_v(a_1, \dots, a_s; q) \text{ and } |V(G)| = n\}.$$

Denote:

$$(1) \quad m(a_1, \dots, a_s) = m = \sum_{i=1}^s (a_i - 1) + 1 \quad \text{and} \quad p = \max\{a_1, \dots, a_s\}.$$

It is easy to prove that $G \xrightarrow{v} (a_1, \dots, a_s) \Rightarrow \chi(G) \geq m$.

If $q \geq m + 1$ then $F_v(a_1, \dots, a_s; q) = m$, since:

- $K_m \xrightarrow{v} (a_1, \dots, a_s)$.

- $K_{m-1} \not\xrightarrow{v} (a_1, \dots, a_s)$.

- For example, $K_5 \xrightarrow{v} (3, 3)$ and $F_v(3, 3; 6) = 5$.

$F_v(a_1, \dots, a_s; m)$:

Theorem 1. (2001, Luczak, Rucinski and Urbanski) *Let a_1, \dots, a_s be positive integers and let m and p be defined by (1). Then:*

(a) $F_v(a_1, \dots, a_s; m) = m + p$.

(b) $K_{m+p} - C_{2p+1} = K_{m-p-1} + \overline{C}_{2p+1}$ is the only graph in $\mathcal{H}_v(a_1, \dots, a_s; m; m + p)$.

· For example, $F_v(3, 3; 5) = 8$ and $K_1 + \overline{C}_7$ is the only 8-vertex graph in $\mathcal{H}_v(3, 3; 5)$.

$F_v(a_1, \dots, a_s; m - 1)$:

- All numbers in the form $F_v(a_1, \dots, a_s; m - 1)$, for which $p \leq 6$ are known.
- Especially interesting is the result $F_v(3, 3; 4) = 14$ (computer aided).
- 2016 Bikov and Nenov: $F_v(2, 2, 7; 8) = 20$. This is the only known number in the form $F_v(a_1, \dots, a_s; m - 1)$ for which $p \geq 7$.
- In the case $p = 7$ we also know the bounds $m + 10 \leq F_v(a_1, \dots, a_s; m - 1) \leq m + 12$.

In this paper we present a new improved algorithm with the help of which we can obtain lower bounds on the numbers $F_v(a_1, \dots, a_s; m - 1)$. Using the new algorithm we improve the lower bound on the numbers $F_v(a_1, \dots, a_s; m - 1)$ when $\max \{a_1, \dots, a_s\} = 7$:

Main Theorem. *Let a_1, \dots, a_s be positive integers, such that $\max \{a_1, \dots, a_s\} = 7$ and $m = \sum_{i=1}^s (a_i - 1) + 1$. Then*

$$F_v(a_1, \dots, a_s; m - 1) \geq m + 11.$$

Bounds on the numbers $F_v(a_1, \dots, a_s; q)$:

Denote:

$\mathcal{S}(m, p)$ is the set of all (b_1, \dots, b_r) (r is not fixed), where b_i are positive integers such that $\max\{b_1, \dots, b_r\} = p$ and $\sum_{i=1}^r (b_i - 1) + 1 = m$.

Let $(a_1, \dots, a_s) \in \mathcal{S}(m, p)$. Then obviously

$$\min_{(b_1, \dots, b_r) \in \mathcal{S}(m, p)} F_v(b_1, \dots, b_r; q) \leq F_v(a_1, \dots, a_s; q) \leq \max_{(b_1, \dots, b_r) \in \mathcal{S}(m, p)} F_v(b_1, \dots, b_r; q).$$

It is easy to prove the following

Proposition 1.

$$\min_{(b_1, \dots, b_r) \in \mathcal{S}(m, p)} F_v(b_1, \dots, b_r; q) = F_v(2_{m-p}, p; q).$$

Theorem 2. (2015, Bikov and Nenov)

(a) *There exists a positive integer r_0 such that*

$$F_v(2_r, p; r + p - 1) = F(2_{r_0}, p; r_0 + p - 1) + r - r_0, \quad r \geq r_0.$$

(b) $r_0 < F_v(2, 2, p; p + 1) - 2p$.

From Theorem 2 it becomes clear, that for fixed p the computation of the members of the infinite sequence $F_v(2_{m-p}, p; m - 1)$, $m \geq p + 2$, is reduced to the computation of its first r_0 members, where $r_0 < F_v(2, 2, p; p + 1) - 2p$. We conjecture that it is enough to know only its first member $F_v(2, 2, p; p + 1)$.

Conjecture 1. *If $p \geq 4$, then*

$$F_v(2_r, p; r + p - 1) = F_v(2, 2, p; p + 1) + r - 2, \quad r \geq 2,$$

or equivalently, if $p \geq 4$, then $r_0 = 2$.

This conjecture is proved for $p = 4, 5$ and 6 . It is also proved that the conjecture is true when $F_v(2, 2, p; p + 1) \leq 2p + 5$. In this paper we prove that Conjecture 1 is also true when $p = 7$. The Main Theorem follows easily from this result.

Algorithms:

- Finding all graphs in $\mathcal{H}(a_1, \dots, a_s; q; n)$ using a brute force approach is practically impossible for $n > 13$.

- We say that G is a maximal graph in $\mathcal{H}(a_1, \dots, a_s; q; n)$ if $G \in \mathcal{H}(a_1, \dots, a_s; q; n)$ but $G + e \notin \mathcal{H}(a_1, \dots, a_s; q; n), \forall e \in E(\overline{G})$. Denote by $\mathcal{H}_{max}(a_1, \dots, a_s; q; n)$ the set of all maximal graphs in $\mathcal{H}(a_1, \dots, a_s; q; n)$.

- Every graph in $\mathcal{H}(a_1, \dots, a_s; q; n)$ can be obtained by removing edges from some graph in $\mathcal{H}_{max}(a_1, \dots, a_s; q; n)$.

Algorithm 1. *The input of the algorithm is the set \mathcal{A} of all graphs in $\mathcal{H}_{max}(a_1 - 1, \dots, a_s; q; n - r)$ with independence number not greater than t . The output of the algorithm is the set \mathcal{B} of all graphs $G \in \mathcal{H}_{max}(a_1, \dots, a_s; q; n)$ with $r \leq \alpha(G) \leq t$.*

With the help of Algorithm 1, finding the maximal graphs in $\mathcal{H}(a_1, \dots, a_s; q; n)$ is reduced to finding maximal graphs with a smaller number of vertices.

· 2016 Bikov and Nenov: $F_v(2, 2, 7; 8) = 20$. Lower bound is proved by showing with the help of Algorithm 1 that $\mathcal{H}_v(2, 2, 7; 8; 19) = \emptyset$.

set	ind. number	maximal graphs	(+ K_7)- graphs
$\mathcal{H}(2, 7; 8; 15)$	≤ 4	1	1
$\mathcal{H}(2, 2, 7; 8; 19)$	$= 4$	0	
$\mathcal{H}(3; 8; 6)$	≤ 3	1	1
$\mathcal{H}(4; 8; 8)$	≤ 3	1	4
$\mathcal{H}(5; 8; 10)$	≤ 3	3	45
$\mathcal{H}(6; 8; 12)$	≤ 3	12	3 104
$\mathcal{H}(7; 8; 14)$	≤ 3	169	4 776 518
$\mathcal{H}(2, 7; 8; 16)$	≤ 3	34	22 896
$\mathcal{H}(2, 2, 7; 8; 19)$	$= 3$	0	
$\mathcal{H}(3; 8; 7)$	≤ 2	1	1
$\mathcal{H}(4; 8; 9)$	≤ 2	1	8
$\mathcal{H}(5; 8; 11)$	≤ 2	3	84
$\mathcal{H}(6; 8; 13)$	≤ 2	10	5 394
$\mathcal{H}(7; 8; 15)$	≤ 2	102	4 984 994
$\mathcal{H}(2, 7; 8; 17)$	≤ 2	2760	380 361 736
$\mathcal{H}(2, 2, 7; 8; 19)$	$= 2$	0	
$\mathcal{H}(2, 2, 7; 8; 19)$		0	

Table 1: Steps in finding all maximal graphs in $\mathcal{H}(2, 2, 7; 8; 19)$

Using Theorem 2 (b), we show that to prove Conjecture 1 in the case $p = 7$ we need to prove the inequalities $F_v(2, 2, 2, 7; 9) > 20$, $F_v(2, 2, 2, 2, 7; 10) > 21$ and $F_v(2, 2, 2, 2, 2, 7; 11) > 22$. We prove these inequalities with the help of the following algorithm, which is a modification of Algorithm 1.

Algorithm 2. *The input of the algorithm are the set \mathcal{A}_1 of all graphs in $\mathcal{H}_{max}(a_1 - 1, \dots, a_s; q; n - r)$ with independence number not greater than t and the set \mathcal{A}_2 of all graphs in $\mathcal{H}_{max}(a_1 - 1, \dots, a_s; q - 1; n - 1)$ with independence number not greater than t . The output of the algorithm is the set \mathcal{B} of all graphs $G \in \mathcal{H}_{max}(a_1, \dots, a_s; q; n)$ with $r \leq \alpha(G) \leq t$.*

Algorithm 2 is based on the fact that if $K_1 + H \in \mathcal{H}(a_1, \dots, a_s; q; n)$, then $H \in \mathcal{H}(a_1 - 1, \dots, a_s; q - 1; n - 1)$. Algorithm 2 uses the graphs already obtained in the proof of the inequality $F_v(2, 2, 7; 8) > 19$ to prove the inequality $F_v(2, 2, 2, 7; 9) > 20$. The proofs of each of the remaining two inequalities uses the graphs obtained in the proof of the previous.

The Main Theorem follows easily from Proposition 1 and Conjecture 1 ($p = 7$).

Let a_1, \dots, a_s be positive integers such that $\max \{a_1, \dots, a_s\} = 7$ and

$$m = \sum_{i=1}^s (a_i - 1) + 1. \text{ Then}$$

$$F_v(a_1, \dots, a_s; m - 1) \geq F_v(2_{m-7}; 7; m - 1) = F_v(2, 2, 7; 8) + m - 9 = m + 11.$$

set	ind. number	max. graphs	max. graphs, no cone v.	$(+K_8)$ - graphs	$(+K_8)$ - graphs, no cone v.
$\mathcal{H}(2, 2, 7; 9; 16)$	≤ 4	1	0	1	0
$\mathcal{H}(2, 2, 2, 7; 9; 20)$	$= 4$	0	0		
$\mathcal{H}(4; 9; 7)$	≤ 3	1	0	1	0
$\mathcal{H}(5; 9; 9)$	≤ 3	1	0	4	0
$\mathcal{H}(6; 9; 11)$	≤ 3	3	0	45	0
$\mathcal{H}(7; 9; 13)$	≤ 3	12	0	3 113	9
$\mathcal{H}(2, 7; 9; 15)$	≤ 3	169	0	4 783 615	7 097
$\mathcal{H}(2, 2, 7; 9; 17)$	≤ 3	36	2	22 918	22
$\mathcal{H}(2, 2, 2, 7; 9; 20)$	$= 3$	0	0		
$\mathcal{H}(4; 9; 8)$	≤ 2	1	0	1	0
$\mathcal{H}(5; 9; 10)$	≤ 2	1	0	8	0
$\mathcal{H}(6; 9; 12)$	≤ 2	3	0	85	1
$\mathcal{H}(7; 9; 14)$	≤ 2	10	0	5 474	80
$\mathcal{H}(2, 7; 9; 16)$	≤ 2	103	1	5 346 982	361 988
$\mathcal{H}(2, 2, 7; 9; 18)$	≤ 2	2845	85	387 948 338	7 586 602
$\mathcal{H}(2, 2, 2, 7; 9; 20)$	$= 2$	0	0		
$\mathcal{H}(2, 2, 2, 7; 9; 20)$		0	0		

Table 2: Steps in finding all maximal graphs in $\mathcal{H}(2, 2, 2, 7; 9; 20)$

set	ind. number	max. graphs	max. graphs, no cone v.	(+ K_9)- graphs	(+ K_9)- graphs, no cone v.
$\mathcal{H}(2, 2, 2, 7; 10; 17)$	≤ 4	1	0	1	0
$\mathcal{H}(2, 2, 2, 2, 7; 10; 21)$	$= 4$	0	0		
$\mathcal{H}(5; 10; 8)$	≤ 3	1	0	1	0
$\mathcal{H}(6; 10; 10)$	≤ 3	1	0	4	0
$\mathcal{H}(7; 10; 12)$	≤ 3	3	0	45	0
$\mathcal{H}(2, 7; 10; 14)$	≤ 3	12	0	3 115	2
$\mathcal{H}(2, 2, 7; 10; 16)$	≤ 3	169	0	4 784 483	868
$\mathcal{H}(2, 2, 2, 7; 10; 18)$	≤ 3	36	0	22 919	1
$\mathcal{H}(2, 2, 2, 2, 7; 10; 21)$	$= 3$	0	0		
$\mathcal{H}(5; 10; 9)$	≤ 2	1	0	1	0
$\mathcal{H}(6; 10; 11)$	≤ 2	1	0	8	0
$\mathcal{H}(7; 10; 13)$	≤ 2	3	0	85	0
$\mathcal{H}(2, 7; 10; 15)$	≤ 2	10	0	5 495	21
$\mathcal{H}(2, 2, 7; 10; 17)$	≤ 2	103	0	5 371 651	24 669
$\mathcal{H}(2, 2, 2, 7; 10; 19)$	≤ 2	2848	3	387 968 658	20 320
$\mathcal{H}(2, 2, 2, 2, 7; 10; 21)$	$= 2$	0	0		
$\mathcal{H}(2, 2, 2, 2, 7; 10; 21)$		0	0		

Table 3: Steps in finding all maximal graphs in $\mathcal{H}(2, 2, 2, 2, 7; 10; 21)$

set	ind. number	max. graphs	max. graphs, no cone v.	(+ K_{10})- graphs	(+ K_{10})- graphs, no cone v.
$\mathcal{H}(2, 2, 2, 2, 7; 11; 18)$	≤ 4	1	0	1	0
$\mathcal{H}(2, 2, 2, 2, 2, 7; 11; 22)$	$= 4$	0	0		
$\mathcal{H}(6; 11; 9)$	≤ 3	1	0	1	0
$\mathcal{H}(7; 11; 11)$	≤ 3	1	0	4	0
$\mathcal{H}(2, 7; 11; 13)$	≤ 3	3	0	45	0
$\mathcal{H}(2, 2, 7; 11; 15)$	≤ 3	12	0	3 116	1
$\mathcal{H}(2, 2, 2, 7; 11; 17)$	≤ 3	169	0	4 784 638	155
$\mathcal{H}(2, 2, 2, 2, 7; 11; 19)$	≤ 3	36	0	22 919	0
$\mathcal{H}(2, 2, 2, 2, 2, 7; 11; 22)$	$= 3$	0	0		
$\mathcal{H}(6; 11; 10)$	≤ 2	1	0	1	0
$\mathcal{H}(7; 11; 12)$	≤ 2	1	0	8	0
$\mathcal{H}(2, 7; 11; 14)$	≤ 2	3	0	85	0
$\mathcal{H}(2, 2, 7; 11; 16)$	≤ 2	10	0	5 502	7
$\mathcal{H}(2, 2, 2, 7; 11; 18)$	≤ 2	103	0	5 374 143	2 492
$\mathcal{H}(2, 2, 2, 2, 7; 11; 20)$	≤ 2	2848	0	387 968 676	18
$\mathcal{H}(2, 2, 2, 2, 2, 7; 11; 22)$	$= 2$	0	0		
$\mathcal{H}(2, 2, 2, 2, 2, 7; 11; 22)$		0	0		

Table 4: Steps in finding all maximal graphs in $\mathcal{H}(2, 2, 2, 2, 2, 7; 11; 22)$

Thank you.